

# *Bratteli-Vershik Models from Branching Rauzy Induction*

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# *Interval Exchange Transformations*

*(In case you arrived late)*



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## Definition (interval exchange transformation)

Let  $\mathcal{A} = \{a_1 < \dots < a_n\}$  be an alphabet of  $n \geq 1$  letters, and fix points  $\ell = z_0 < z_1 < \dots < z_n = r$ . Consider the partition of  $I = [\ell, r)$  into the subintervals  $I_{a_i} = [z_{i-1}, z_i)$ ,  $1 \leq i \leq n$ . Given  $\pi \in S_{\mathcal{A}}$ , the associated  $n$ -interval exchange transformation is the map  $T : I \rightarrow I$  acting on each sub-interval by a translation

$$T(z) = z + \tau_a, \quad z \in I_a, \quad \tau_a = \sum_{\pi^{-1}(b) < \pi^{-1}(a)} |I_b| - \sum_{b < a} |I_b|.$$

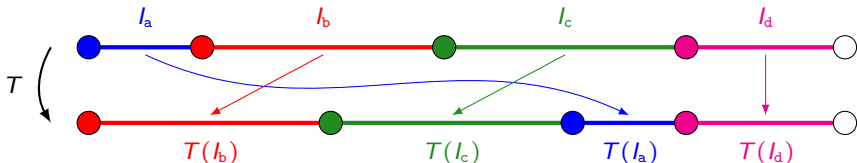
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# *Formal Discontinuities*

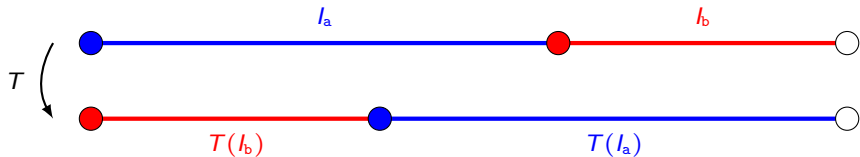
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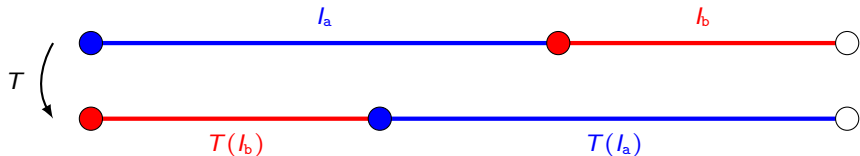
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### Definition (Formal discontinuity points)

Let  $(I, T)$  be an interval exchange transformation with a partition  $\mathcal{P} = \{I_{a_1}, \dots, I_{a_n}\}$ , where  $I_{a_i} = [z_{i-1}, z_i)$  and  $\ell = z_0 < z_1 < \dots < z_n = r$ . The set of *formal discontinuity points* of  $T$ , denoted by  $\mathcal{D}(T)$ , is the set of points  $z_1, \dots, z_{n-1}$ .

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All regular interval exchanges are minimal, but the converse is false.



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$$X = ([\ell, r] \setminus \mathcal{O}(T)) \cup \{z^-, z^+ \mid z \in \mathcal{O}(T)\},$$

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## Definition (Cantor version of an interval exchange)

Let  $(I, T)$  be a minimal interval exchange with  $X$  as defined previously. Define the map  $\phi : X \rightarrow X$  by  $\phi(x) = T(x)$  for  $x \in X \setminus \{z^-, z^+ \mid z \in \mathcal{O}(T)\}$ ,  $\phi(x^-) = (T(x))^-$ ,  $\phi(x^+) = (T(x))^+$  for  $x \in \mathcal{O}(T)$ .

The dynamical system  $(X, \phi)$  is called the *Cantor version of  $T$* .

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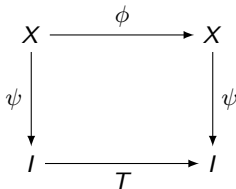
Let  $(I, T)$  be a minimal interval exchange and let  $(X, \phi)$  be its Cantor version. Define the map  $\psi : X \rightarrow I$  by  $\psi(x^-) = \psi(x^+) = x$  for  $x \in \mathcal{O}(T)$  and  $\psi(x) = x$  otherwise. The map  $\psi$  is a surjection. It is also an injection off the duplicated points of  $X$ .



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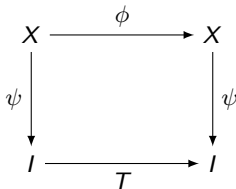
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Similarly, there is an inclusion map  $\iota : I \rightarrow X$  defined by  $\iota(x) = x$  when  $x \in I \setminus \mathcal{O}(T)$ , and  $\iota(x) = x^+$  when  $x \in \mathcal{O}(T)$  such that  $(\iota \circ T)(x) = (\phi \circ \iota)(x)$  for all  $x \in I$ .

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Let  $(I, T)$  be a regular interval exchange and let  $(X, \phi)$  be its Cantor version. If  $D = [a, b)$  is a (branching) Rauzy induction domain for  $T$ , then the corresponding clopen set

$$D_X = [a^+, b^-] \subset X$$

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Thus, Cantorization and induction commute.



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### Definition (Partition in towers)

Let  $(X, T)$  be an invertible Cantor minimal system. Let  $\{I_1, \dots, I_n\}$  be a collection of pairwise disjoint subsets of  $X$ , and let  $\{h_1, \dots, h_n\}$  be a set of  $n$  positive integers. If the collection  $\{T^i(I_j) \mid 0 \leq i \leq h_j, 1 \leq j \leq n\}$  is a partition of  $X$ , it is called a *partition in towers* of  $(X, T)$ .

The intervals  $I_1, \dots, I_n$  are called tower bases, the integers  $h_1, \dots, h_n$  are called tower heights, and the set  $\{T^i(I_j) \mid 0 \leq i \leq h_j\}$  is called the  $j^{\text{th}}$  tower.

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The union  $I_1 \cup \dots \cup I_n$  is called the *basis* of the partition in towers.

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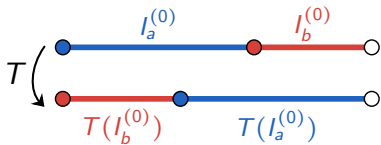
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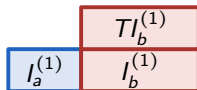
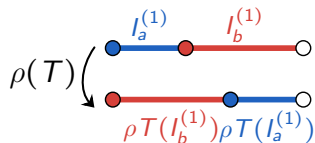
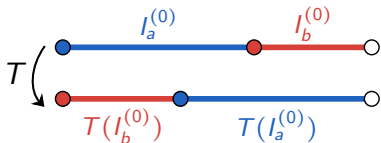
A sequence of partitions in towers  $(\mathcal{B})_{n \geq 0}$  is said to be nested if  $\mathcal{B}_{n+1}$  is nested in  $\mathcal{B}_n$  for all  $n \geq 0$ .

*Partitions in Towers from the Fibonacci IET*

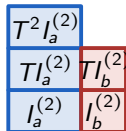
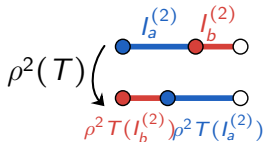
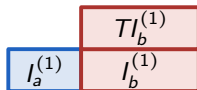
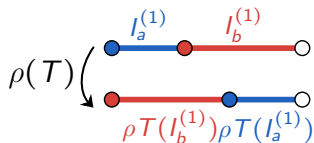
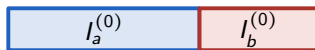
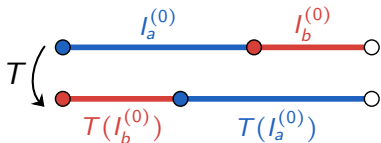
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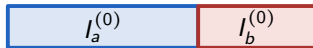


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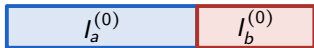
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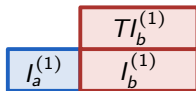
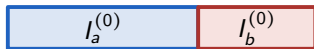
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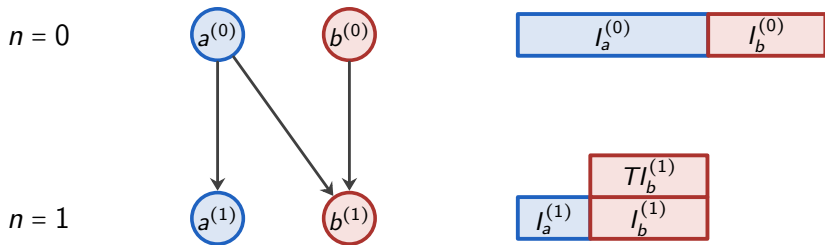


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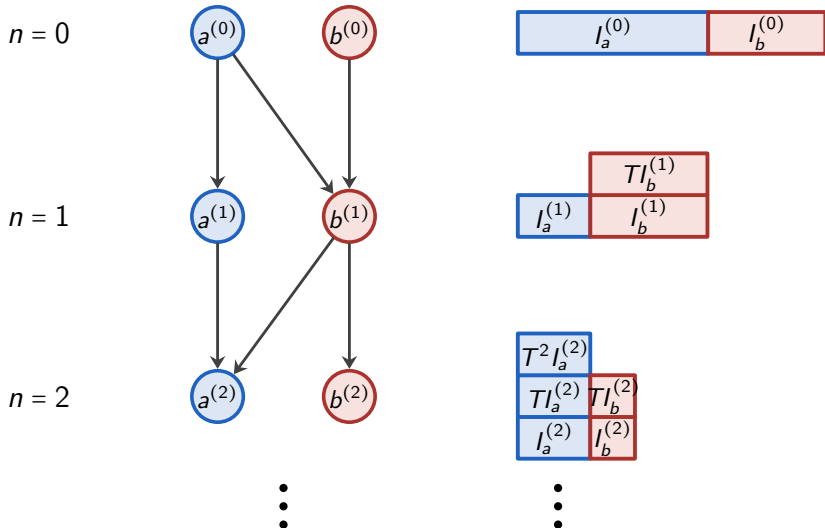
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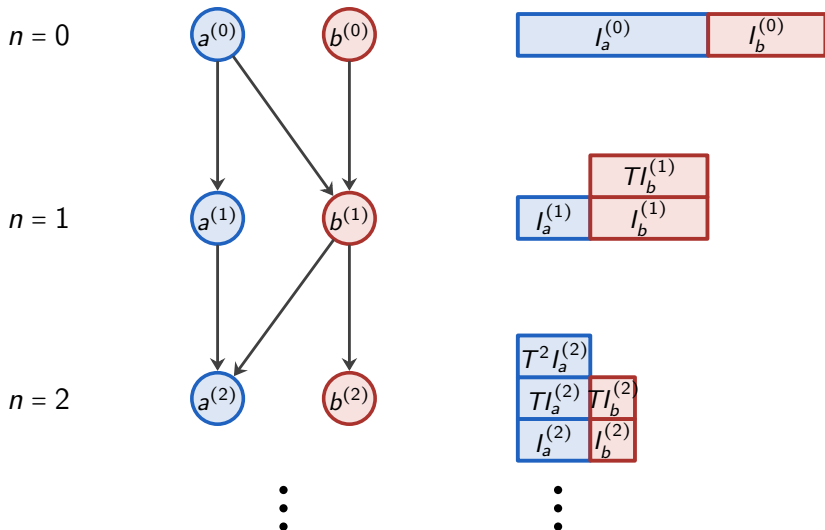
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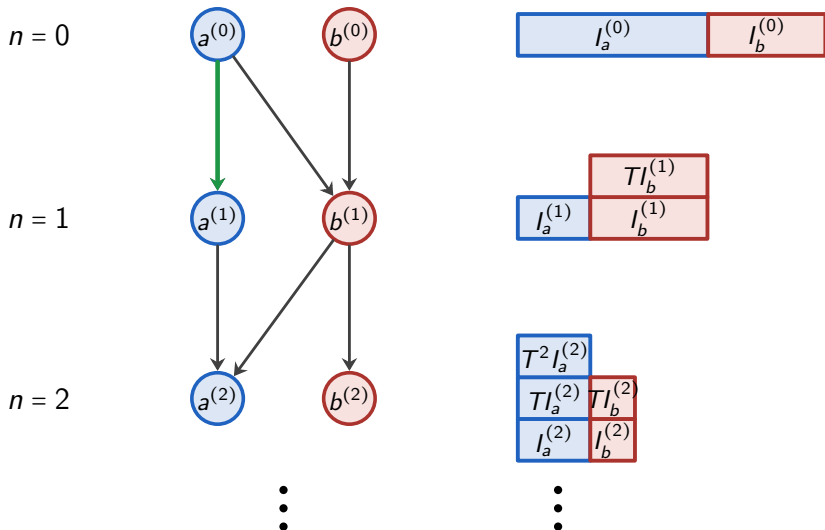
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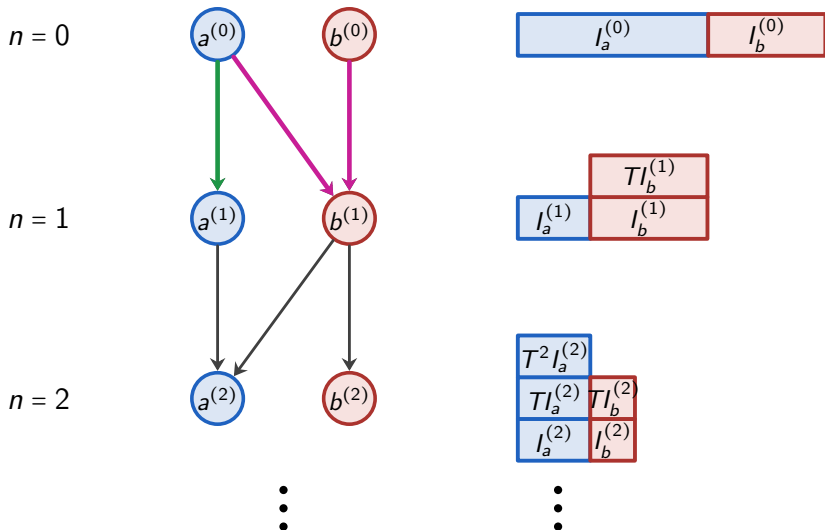
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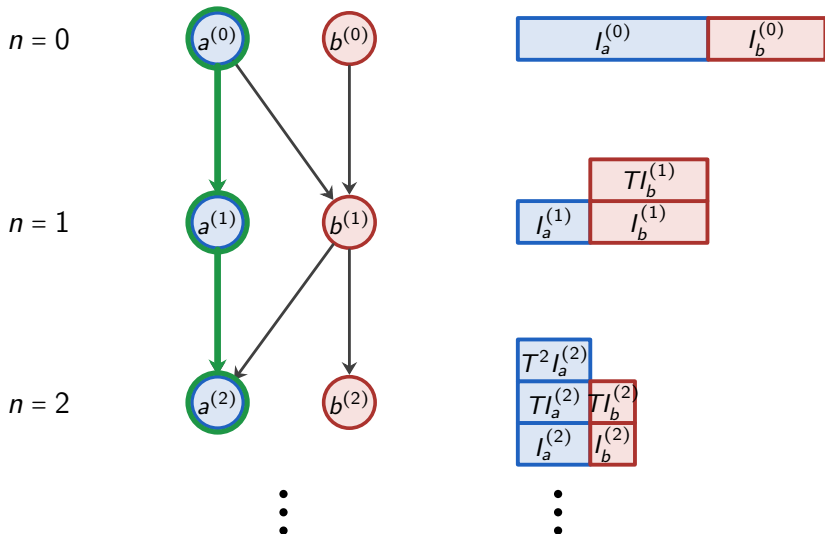
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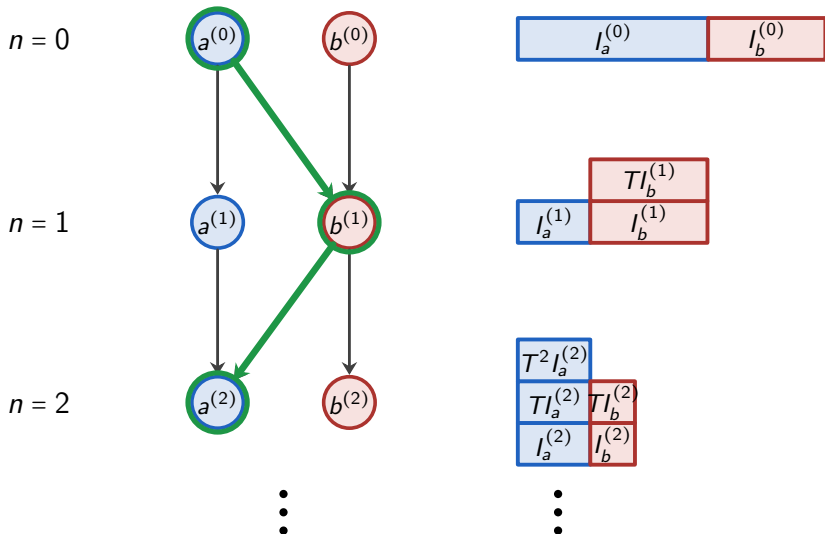
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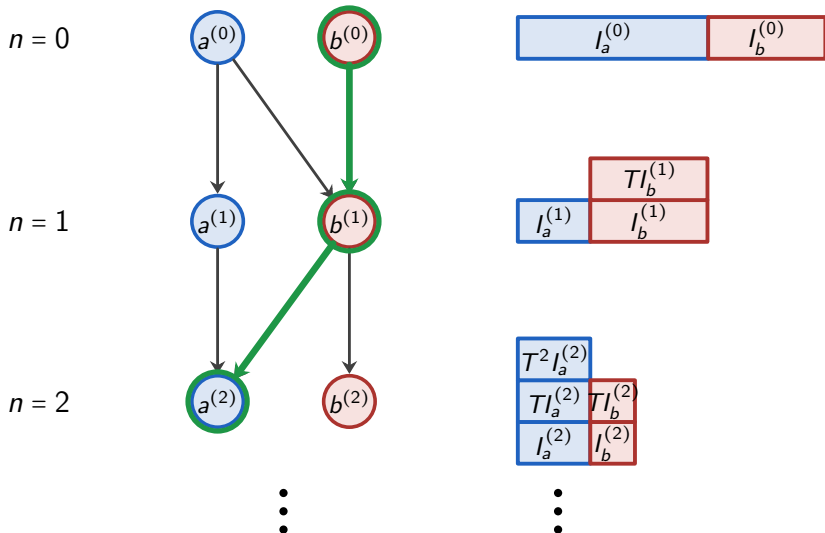
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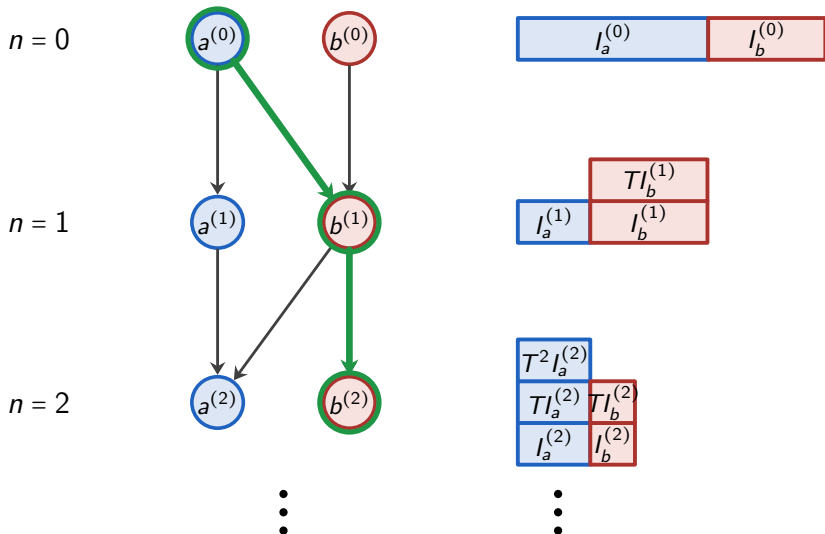
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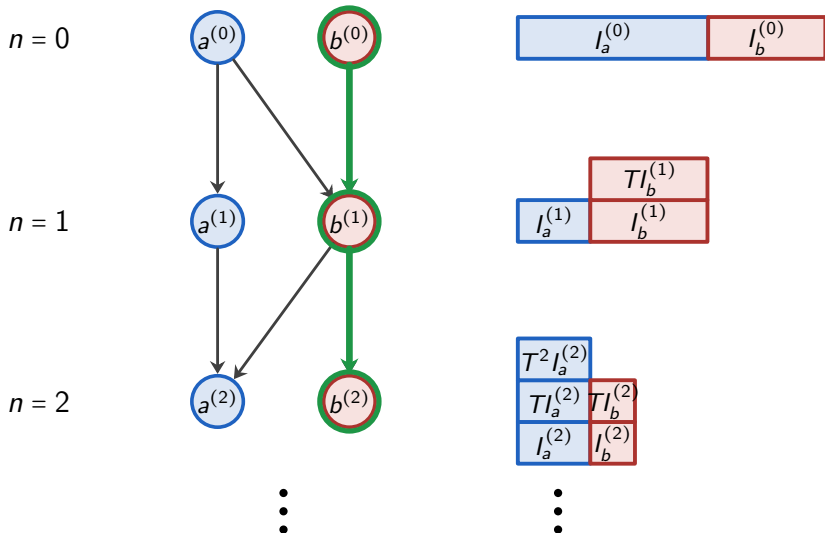
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### Definition (Bratteli diagram)

A *Bratteli diagram* is an infinite directed graph  $\mathcal{B} = (V, E)$  whose vertex set can be partitioned into levels  $V_0, V_1, V_2, \dots$  and whose edges  $E$  can be similarly partitioned into sets  $E_1, E_2, E_3, \dots$  such that edges in  $E_i$  only go from vertices in  $V_{i-1}$  to vertices in  $V_i$ , and such that  $V_0$  is a one-vertex set called the *hat* of  $\mathcal{B}$ .

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Steps to construct a Bratteli diagram from a nested sequence of partitions in towers :

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A *Bratteli diagram* is an infinite directed graph  $\mathcal{B} = (V, E)$  whose vertex set can be partitioned into levels  $V_0, V_1, V_2, \dots$  and whose edges  $E$  can be similarly partitioned into sets  $E_1, E_2, E_3, \dots$  such that edges in  $E_i$  only go from vertices in  $V_{i-1}$  to vertices in  $V_i$ , and such that  $V_0$  is a one-vertex set called the *hat* of  $\mathcal{B}$ .

It is called *ordered* if there is also a partial order  $\leq$  on  $E$  such that two edges in  $E$  are comparable if and only if they have the same range vertex.

Steps to construct a Bratteli diagram from a nested sequence of partitions in towers :

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Thus, nested tower partitions naturally produce ordered Bratteli diagrams.

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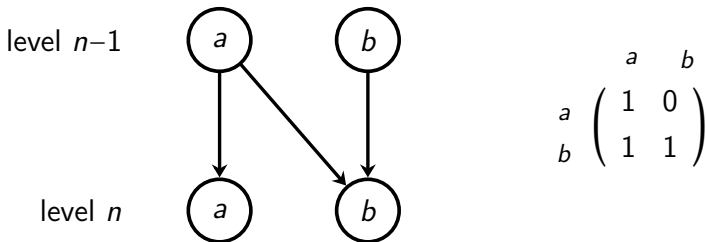
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Products of incidence matrices count paths across several levels.

In particular, if sufficiently long products have all entries positive, then the diagram is *simple*.

## The Incidence Matrix of One Level



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Let  $\mathcal{B} = (V, E)$  be a Bratteli diagram. The *infinite path space* of  $\mathcal{B}$  is the set

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## Definition (Path cylinder set)

A path cylinder set is a set  $[e_1, e_2, \dots, e_n]_E = \{(f_1, f_2, \dots) \in X_E \mid f_i = e_i, 1 \leq i \leq n\}$ .

# *Maximal and Minimal Paths*

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### Definition (Maximal/minimal path)

Let  $X_E$  be the infinite path space for an ordered Bratteli diagram and let  $(e_1, e_2, \dots, e_i, \dots) \in X_E$ . The path  $(e_1, e_2, \dots, e_i, \dots)$  is said to be *maximal* (resp. *minimal*) if each  $e_i$  is maximal (resp. minimal) on level  $i$ . We denote the set of such paths by  $X_E^{\max}$  (resp.  $X_E^{\min}$ ).

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### Definition (Properly ordered Bratteli diagram)

Given an ordered Bratteli diagram  $\mathcal{B} = (V, E, \leq)$  and its associated infinite path space,  $X_E$ ,  $\mathcal{B}$  is said to be *properly ordered* if it is simple and the sets  $X_E^{\max}$  and  $X_E^{\min}$  are singletons.

# *The Vershik Map*

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### Definition (Successor path in reverse lexicographic order)

Let  $\mathcal{B} = (V, E, \leq)$  be an ordered Bratteli diagram, and let  $(e_1, e_2, \dots, e_i, \dots) \in X_E$ , where  $e_i$  is the first non-maximal edge in the path. The *successor path* of  $(e_1, e_2, \dots, e_i, \dots)$  is the path  $(f_1, f_2, \dots, f_{i-1}, f_i, e_{i+1}, \dots)$  obtained by setting  $e_i$  to its successor,  $f_i$ , on level  $i$ , then taking each  $f_{i-1}$  to be minimal (up to level-wise source-range agreement) for levels  $1 \leq j \leq i-1$ . The successor is undefined if the path is in  $X_E^{\max}$ .

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## Definition (Vershik map)

Let  $X_E$  be the path space of a properly ordered Bratteli diagram  $\mathcal{B}$  and let  $x \in X_E$ . The *Vershik map*  $\varphi_E : X_E \rightarrow X_E$  associated with  $\mathcal{B}$  is the map defined by  $\varphi(x) = \text{successor}(x)$  if  $x \notin X_E^{\max}$ , and  $\varphi(x) = x_{\min}$  if  $x \in X_E^{\max}$ , where  $x_{\min}$  is the single path in  $X_E^{\min}$ .

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Can we relate the behavior of the Vershik map on the Bratteli diagram to the dynamical system from which the Bratteli diagram was generated?

# *Refining Sequences of Partitions in Towers*

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## Definition (Refining sequence of partitions in towers)

A sequence  $(\mathfrak{P}_n)_{n \geq 0}$  of partitions in towers for a system  $(X, T)$  is said to be *refining* if

1. For the sequence of bases  $(B_n)_{n \geq 0}$  of  $(\mathfrak{P}_n)_{n \geq 0}$ ,  $\bigcap_{n \geq 0} B_n$  is a single point.
2. The sequence  $(\mathfrak{P}_n)_{n \geq 0}$  is nested.
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## Theorem (Herman, Putnam, Skau)

Let  $(\mathfrak{P}_n)_{n \geq 0}$  be a refining sequence of partitions in towers for a minimal invertible Cantor system  $(X, T)$ . Then  $(X, T) \cong (X_E, \varphi_E)$ , where  $X_E$  is the path space of the ordered Bratteli diagram arising from  $(\mathfrak{P}_n)_{n \geq 0}$ .

# *The Gjerde-Johansen Construction*

## Theorem (Gjerde, Johansen (2002))

Let  $T$  be a minimal IET and let  $\phi$  be its Cantor version. There exists a simple ordered Diagram (obtained via an acceleration of right cuts) upon which the Vershik map is conjugate to  $\phi$ .

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### Proposition (Dolce, H. 2026+)

Let  $(I, T)$  be a regular IET and  $(X, \phi)$  and  $\chi \in \{\lambda, \rho\}^\omega$ . The sequence of partitions in towers associated with  $\chi$  and  $\phi$  is nesting and generates the topology of  $X$ .

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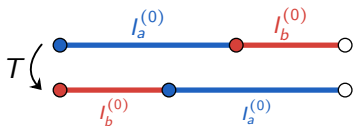
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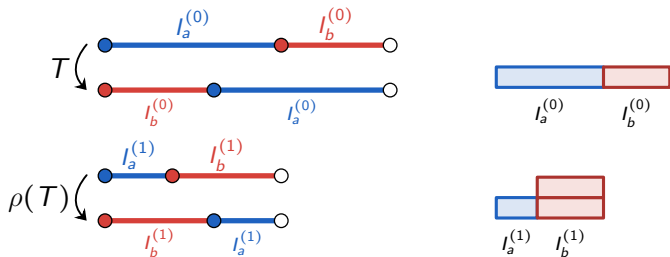
Let  $(I, T)$  and  $(X, \phi)$  be as above. There exist countably many paths in  $\chi \in \{\lambda, \rho\}^\omega$  for which the sequence of bases of the associated partitions in towers has an intersection with cardinality 2.

# *Fibonacci Revisited*

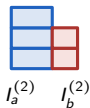
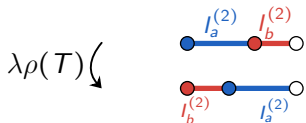
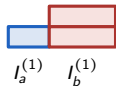
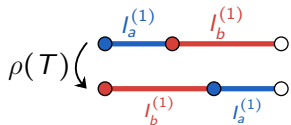
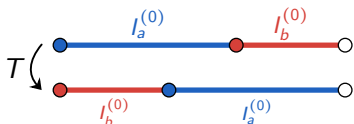
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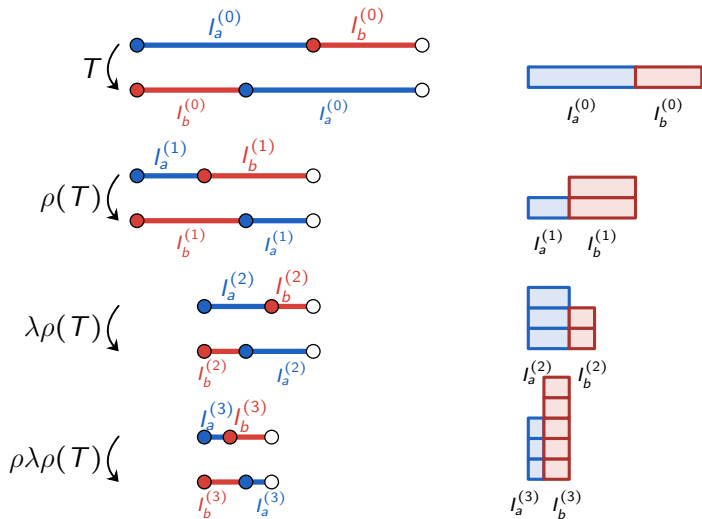
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# *Obtaining a Conjugate Map Anyway*

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### Theorem (Dolce, H. 2026+)

Let  $(I, T)$  be a regular interval exchange and  $(X, \phi)$  be its Cantorized version. For each  $\chi \in \{\lambda, \rho\}^\omega$ , there exists a homeomorphism  $\phi_E$  on the path space  $X_E$  of the Bratteli diagram associated with  $\phi$  and  $\chi$  such that  $(X_E, \phi_E) \cong (X, \phi)$ . Furthermore,  $\phi_E$  agrees with the successor map on paths  $X_E \setminus X_E^{\max}$ .

**Merci Beaucoup !**